

A NEW DECISION PROCEDURE FOR FINITE SETS AND CARDINALITY IN SMT

K. Bansal* and C. Barrett* and A. Reynolds+ and C. Tinelli+

June 30, 2016

*New York University

+The University of Iowa

New tableaux-style calculus for many-sorted quantifier-free
theory of finite sets with cardinality constraints

New tableaux-style calculus for **many-sorted quantifier-free**
theory of **finite sets** with **cardinality** constraints

Motivation: formalizes modular, incremental T -solver that can be integrated into SMT solvers

New tableaux-style calculus for **many-sorted quantifier-free**
theory of **finite sets** with **cardinality** constraints

Motivation: formalizes modular, incremental *T*-solver that can be integrated into SMT solvers

Target use: static analysis tools (Leon, LiquidHaskell, ...)

New tableaux-style calculus for **many-sorted quantifier-free**
theory of **finite sets** with **cardinality** constraints

Motivation: formalizes modular, incremental *T*-solver that can be integrated into SMT solvers

Target use: static analysis tools (Leon, LiquidHaskell, ...)

Current restriction: Element sort must be infinite

New tableaux-style calculus for **many-sorted quantifier-free**
theory of **finite sets** with **cardinality** constraints

Motivation: formalizes modular, incremental *T*-solver that can be integrated into SMT solvers

Target use: static analysis tools (Leon, LiquidHaskell, ...)

Current restriction: Element sort must be infinite

Implementation: initial version, fully integrated in CVC4

(Cantone and Zarba, FTP 1998)

Finite sets with $\{_ \}$, \cup , \cap , \setminus , \in , \subseteq . Tableaux-based procedure.

(Zarba, FroCoS 2002)

As above plus $|_ |$ and LIA constraints. Theoretical results.
Highly inefficient in practice.

(Suter et al., VMCAI 2011)

Boolean Algebras and Presburger Arithmetic, no $\{_ \}$ and \in .
Based on Venn-region computation. Relatively inefficient in practice.

- A more **efficient, decision procedure** for sets of literals
- **incremental** and **modular**
- Capitalizes on **separate** solvers for CC and for LIA
- Reasons **modulo equality** (over elements and over sets)
- Tries to **minimize** number of **Venn regions** that need to be considered
- Introduces Venn regions **lazily**

Try to build a **model** for input constraint set

Try to build a **model** for input constraint set

1. Build (partial) candidate model while ignoring cardinality constraints
 - For each element term e and *relevant* set term s , determine if e is in t or not
 - Work modulo equality
 - Propagate equalities

Try to build a **model** for input constraint set

1. Build (partial) candidate model while ignoring cardinality constraints
 - For each element term e and *relevant* set term s , determine if e is in t or not
 - Work modulo equality
 - Propagate equalities
2. Adjust candidate model to accommodate cardinality constraints
 - Decompose each set into a **partition** of non-empty **regions**
 - **Propagate** cardinality constraints over these decompositions

$\{_ \} : \text{Elem} \rightarrow \text{Set}$ $\sqcup, \sqcap, \setminus : \text{Set} \times \text{Set} \rightarrow \text{Set}$ $\emptyset : \text{Set}$ $\sqsubseteq : \text{Set} \times \text{Set}$ $\in : \text{Elem} \times \text{Set}$ $|_ | : \text{Set} \rightarrow \text{Card}$ $n : \text{Card}$ for all $n \in \mathbb{N}$ $- : \text{Card} \rightarrow \text{Card}$ $+ : \text{Card} \times \text{Card} \rightarrow \text{Card}$ $< : \text{Card} \times \text{Card}$ $\geq : \text{Card} \times \text{Card}$ $\approx : \alpha \times \alpha$ for $\alpha \in \{\text{Set}, \text{Elem}, \text{Card}\}$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in s or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

x	y	
		A
		B
		C
		D, E
		F
		$A \sqcup B, C \cap D$
		$E \sqcup F$
		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in s or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$
2. $D \approx E$
- 3. $x \in C$
4. $x \in E \sqcup F$
5. $x \notin D$
6. $y \notin C \cap D$
7. $y \in A \sqcup D$

x	y	
		A
		B
✓		C
		D, E
		F
		$A \sqcup B, C \cap D$
		$E \sqcup F$
		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

→ 4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

x	y	
		A
		B
✓		C
		D, E
		F
		$A \sqcup B, C \cap D$
✓		$E \sqcup F$
		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in s or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

→ 5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

x	y	
		A
		B
✓		C
✗		D, E
		F
		$A \sqcup B, C \cap D$
✓		$E \sqcup F$
		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in s or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

x	y	
		A
		B
	✓	C
→	✗	D, E
	✓	F
		$A \sqcup B, C \cap D$
→	✓	$E \sqcup F$
		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in s or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

x	y	
		A
		B
	✓	C
→	✗	D, E
	✓	F
	✗	$A \sqcup B, C \cap D$
	✓	$E \sqcup F$
		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

x	y	
X		A
X		B
✓		C
X		D, E
✓		F
→ X		$A \sqcup B, C \cap D$
✓		$E \sqcup F$
		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in s or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

	x	y	
→	X		A
	X		B
	✓		C
→	X		D, E
	✓		F
	X		$A \sqcup B, C \cap D$
	✓		$E \sqcup F$
	X		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

→ 6. $y \notin C \cap D$

7. $y \in A \sqcup D$

x	y	
X		A
X		B
✓		C
X		D, E
✓		F
X	X	$A \sqcup B, C \cap D$
✓		$E \sqcup F$
X		$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

➔ 7. $y \in A \sqcup D$

x	y	
✗		A
✗		B
✓		C
✗		D, E
✓		F
✗	✗	$A \sqcup B, C \cap D$
✓		$E \sqcup F$
✗	✓	$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

	x	y	
	X	X	A
	X	X	B
	✓		C
	X		D, E
	✓		F
→	X	X	$A \sqcup B, C \cap D$
	✓		$E \sqcup F$
	X	✓	$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

	x	y	
→	X	X	A
	X	X	B
	✓		C
	X	✓	D, E
	✓		F
	X	X	$A \sqcup B, C \cap D$
	✓		$E \sqcup F$
→	X	✓	$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

	x	y	
	X	X	A
	X	X	B
	✓		C
→	X	✓	D, E
	✓		F
	X	X	$A \sqcup B, C \cap D$
	✓	✓	$E \sqcup F$
	X	✓	$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

	x	y	
	X	X	A
	X	X	B
	✓	X	C
→	X	✓	D, E
	✓		F
→	X	X	$A \sqcup B, C \cap D$
	✓	✓	$E \sqcup F$
	X	✓	$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 1

1. $A \sqcup B \approx C \cap D$

2. $D \approx E$

3. $x \in C$

4. $x \in E \sqcup F$

5. $x \notin D$

6. $y \notin C \cap D$

7. $y \in A \sqcup D$

x	y	
X	X	A
X	X	B
✓	X	C
X	✓	D, E
✓		F
X	X	$A \sqcup B, C \cap D$
✓	✓	$E \sqcup F$
X	✓	$A \sqcup D$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 2

1. $x \notin A \sqcup B$
2. $x \in A \cap C$

x	
	A
	B
	C
\times	$A \sqcup B$
\checkmark	$A \cap C$

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 2

1. $x \notin A \sqcup B$
2. $x \in A \sqcap C$

x	
x	A
x	B
	C
\rightarrow x	$A \sqcup B$
\checkmark	$A \sqcap C$

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 2

1. $x \notin A \sqcup B$
2. $x \in A \sqcap C$

x	
x ✓	A
x	B
✓	C
x	$A \sqcup B$
→ ✓	$A \sqcap C$

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 2

1. $x \notin A \sqcup B$
2. $x \in A \cap C$

x	
x ✓	A
x	B
✓	C
x	$A \sqcup B$
✓	$A \cap C$

Contradiction!

PROCESSING SET CONSTRAINTS

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 3

...

$x \approx y$

...

x	y	
	✓	A
✗		B
✓		C
	✗	D
...

- For each element term e and *relevant* set term s , determine if e is in t or not
- Work modulo equality

Example 3

→ ...
 $x \approx y$
...

x, y	
✓	A
✗	B
✓	C
✗	D
...	...

- Propagate equalities

Example 4

...

$$A \approx \{y\}$$

$$A \approx B$$

$$x \in A$$

$$x \in A \sqcup C$$

$$y \in B \sqcup C$$

...

x	y	
✓	✓	A, B, {y}
...	...	A \sqcup C
...	...	B \sqcup C
...

- Propagate equalities

Example 4

...

$$A \approx \{y\}$$

$$A \approx B$$

→ $x \in A$

$$x \in A \sqcup C$$

$$y \in B \sqcup C$$

...

x, y	
✓	$A, B, \{y\}$
...	$A \sqcup C$
...	$B \sqcup C$
...	...

- Propagate equalities

Example 4

...

$$A \approx \{y\}$$

→ $A \approx B$

$$x \in A$$

$$x \in A \sqcup C$$

$$y \in B \sqcup C$$

...

x, y	
✓	$A, B, \{y\}$
...	$A \sqcup C, B \sqcup C$
...	...

Simplifying requirements:

- Only variables for elements
- No applications of \sqsubseteq
- All constraints in *flat form* ($x \in A \text{ op } B, A \text{ op } B \approx C \text{ op } D$)
- Each set variable occurs **at most once** in non-variable terms

Simplifying requirements:

- Only variables for elements
- No applications of \sqsubseteq
- All constraints in *flat form* ($x \in A \text{ op } B, A \text{ op } B \approx C \text{ op } D$)
- Each set variable occurs **at most once** in non-variable terms

No loss of generality

Cardinality graph $G = (N, E)$

N , **selected** terms in or constructed from input problem

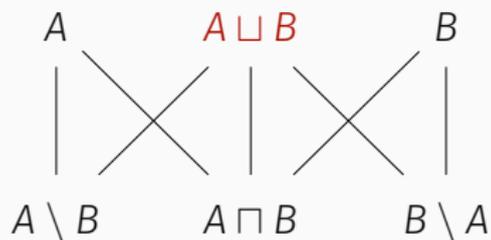
E , s.t. the leaves of each subtree form a **partition** of its root

Cardinality graph $G = (N, E)$

N , **selected** terms in or constructed from input problem

E , s.t. the leaves of each subtree form a **partition** of its root

Example



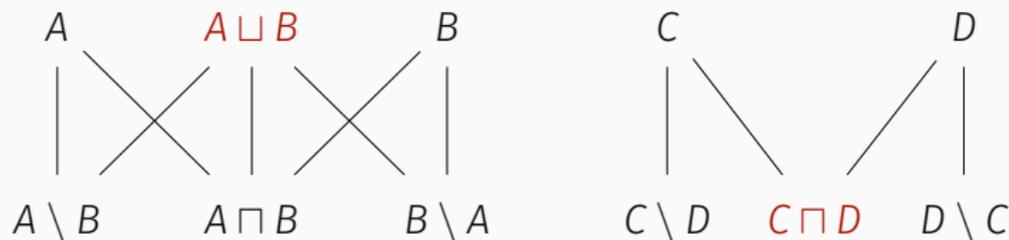
Note: $|A \sqcup B| = |A \setminus B| + |A \cap B| + |B \setminus A|$ $|A| = |A \setminus B| + |A \cap B|$

Cardinality graph $G = (N, E)$

N , **selected** terms in or constructed from input problem

E , s.t. the leaves of each subtree form a **partition** of its root

Example



Note: $|A \sqcup B| = |A \setminus B| + |A \cap B| + |B \setminus A|$ $|A| = |A \setminus B| + |A \cap B|$

$$x \notin D$$

$$y \notin A \sqcup B$$

$$x \in E \sqcup F$$

$$x \in C$$

$$y \in D$$

$$A \sqcup B \approx C \cap D$$

$$D \approx E$$

$$|A \sqcup B| \geq 4$$

$$|C| + |D| \leq 10$$

$$|E \sqcup F| \geq 100$$

$$x \notin D$$

$$y \notin A \sqcup B$$

$$x \in E \sqcup F$$

$$x \in C$$

$$y \in D$$

$$A \sqcup B \approx C \cap D$$

$$D \approx E$$

$$|A \sqcup B| \geq 4$$

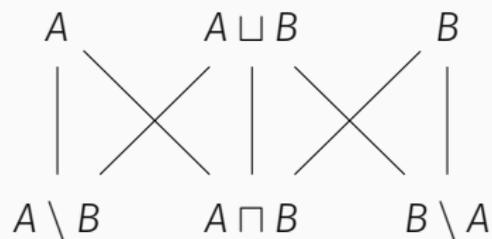
$$|C| + |D| \leq 10$$

$$|E \sqcup F| \geq 100$$

PROCESSING CARDINALITY CONSTRAINTS: EXAMPLE

$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

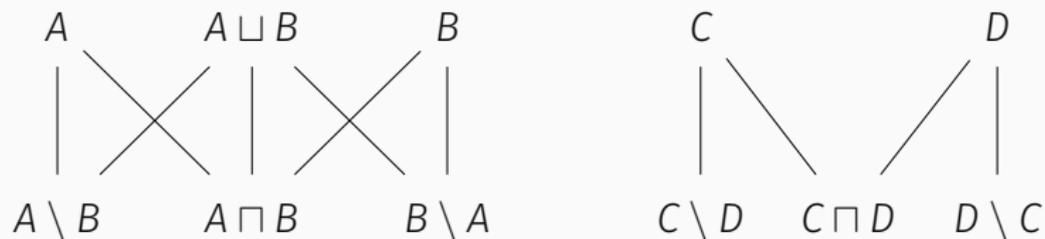
$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 10, |E \sqcup F| \geq 100$



PROCESSING CARDINALITY CONSTRAINTS: EXAMPLE

$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

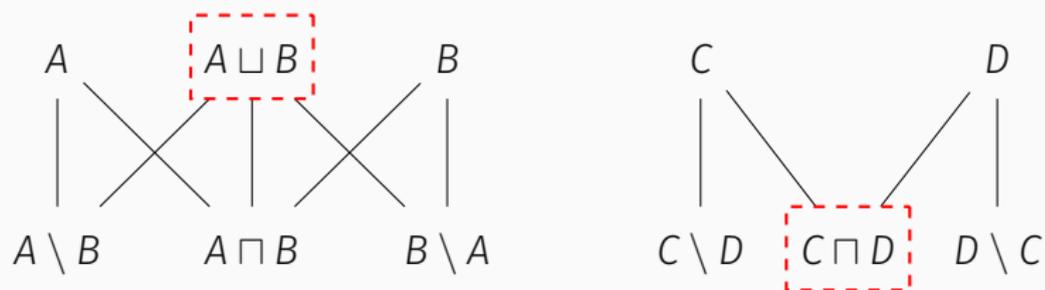
$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 10, |E \sqcup F| \geq 100$



PROCESSING CARDINALITY CONSTRAINTS: EXAMPLE

$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 10, |E \sqcup F| \geq 100$

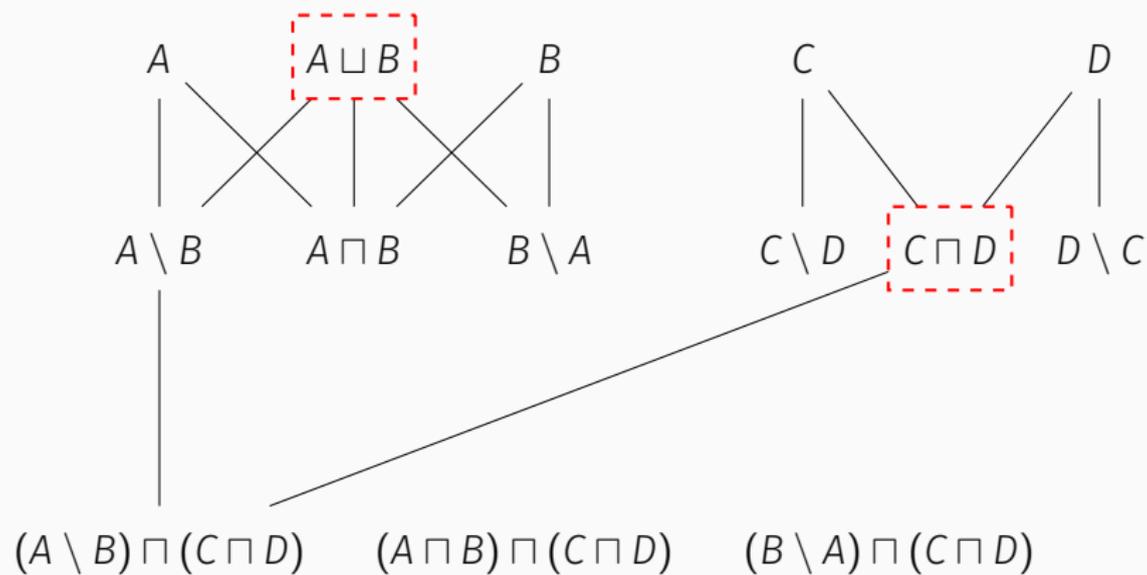


$(A \setminus B) \cap (C \cap D)$ $(A \cap B) \cap (C \cap D)$ $(B \setminus A) \cap (C \cap D)$

PROCESSING CARDINALITY CONSTRAINTS: EXAMPLE

$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

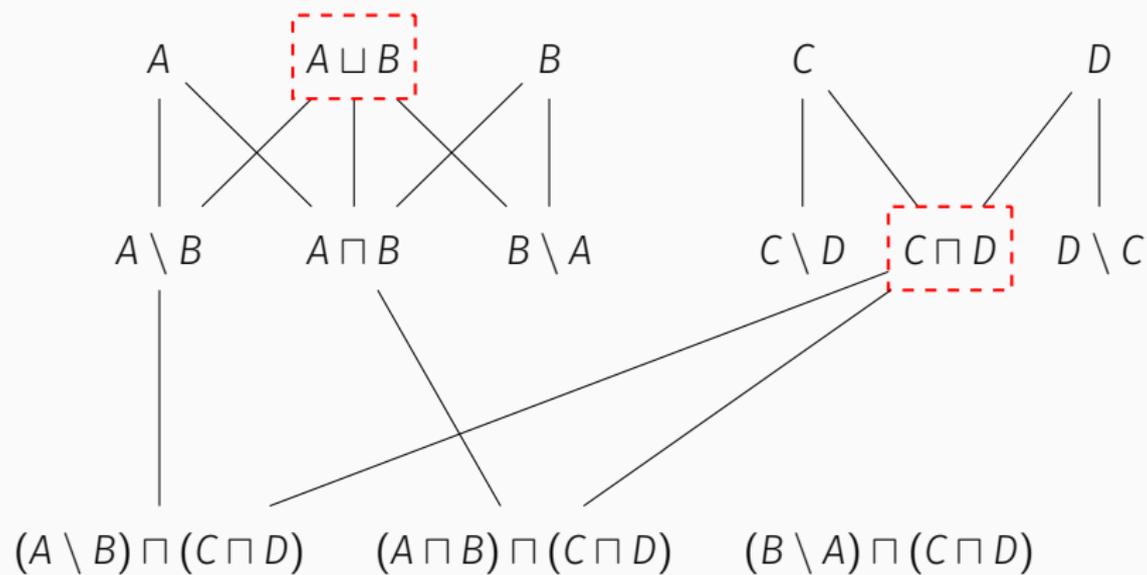
$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 10, |E \sqcup F| \geq 100$



PROCESSING CARDINALITY CONSTRAINTS: EXAMPLE

$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

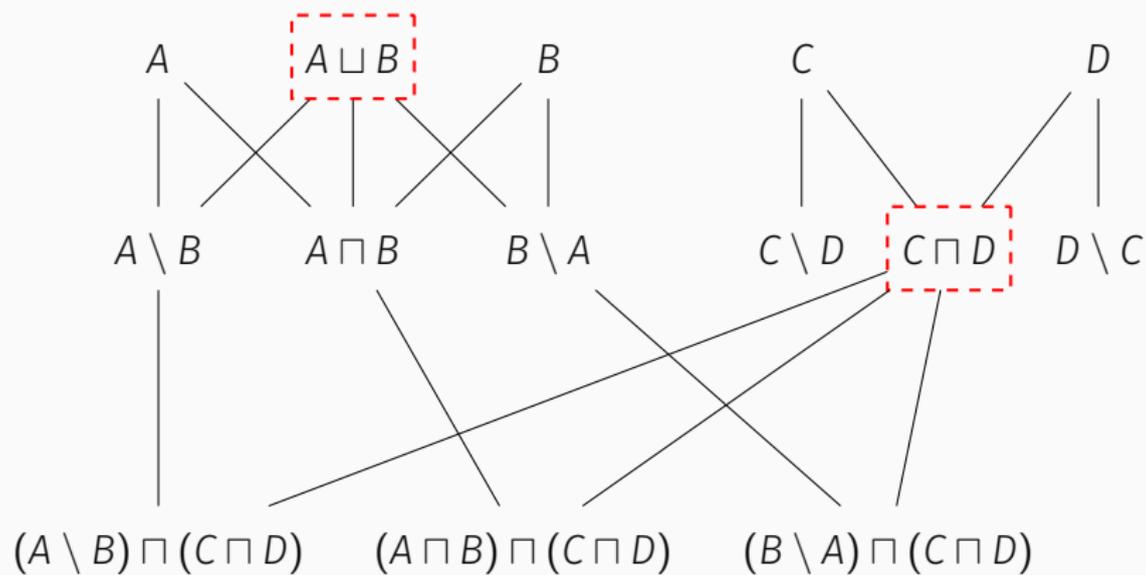
$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 10, |E \sqcup F| \geq 100$



PROCESSING CARDINALITY CONSTRAINTS: EXAMPLE

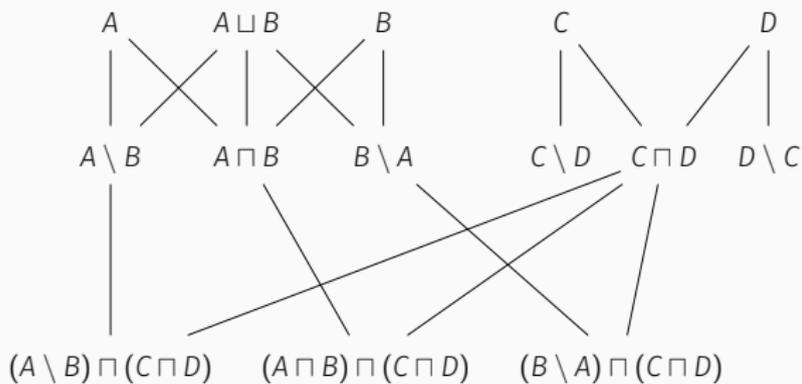
$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 10, |E \sqcup F| \geq 100$



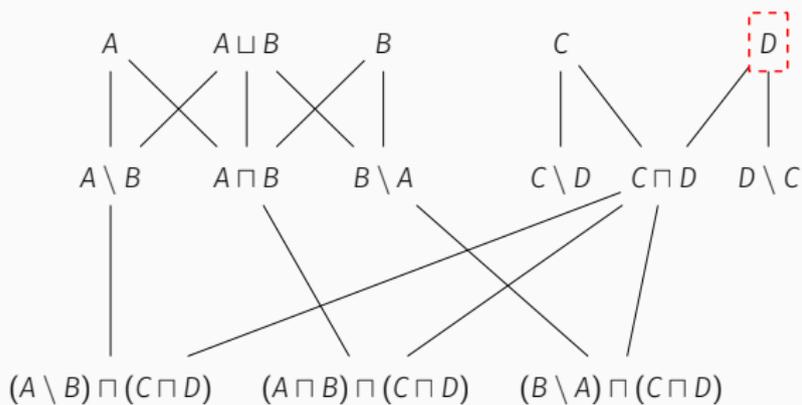
$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 9, |E \sqcup F| \geq 100$



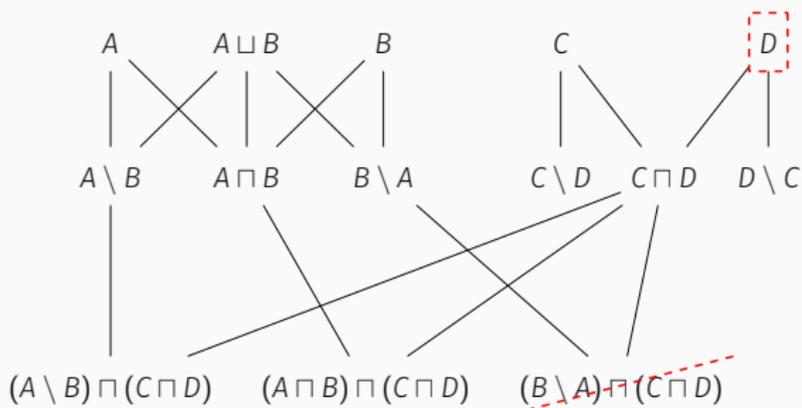
$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 9, |E \sqcup F| \geq 100$



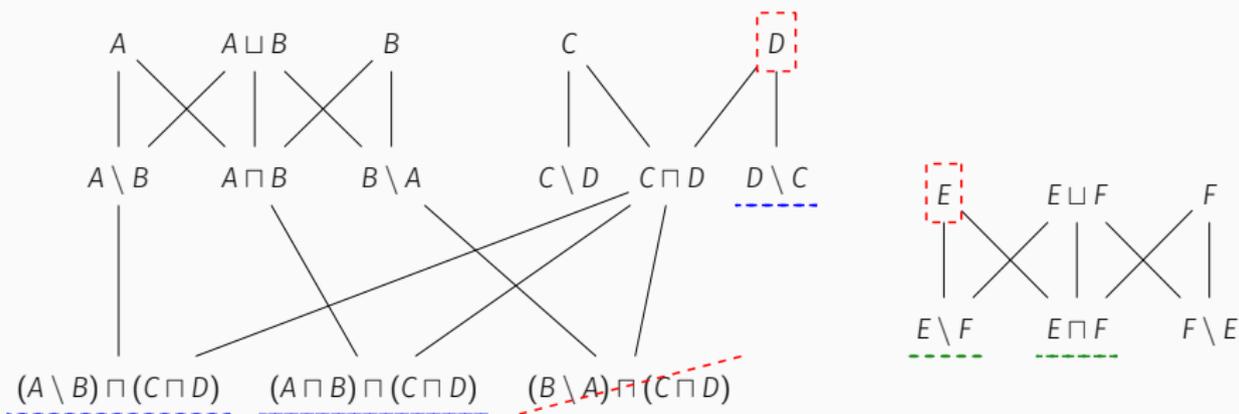
$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 9, |E \sqcup F| \geq 100$



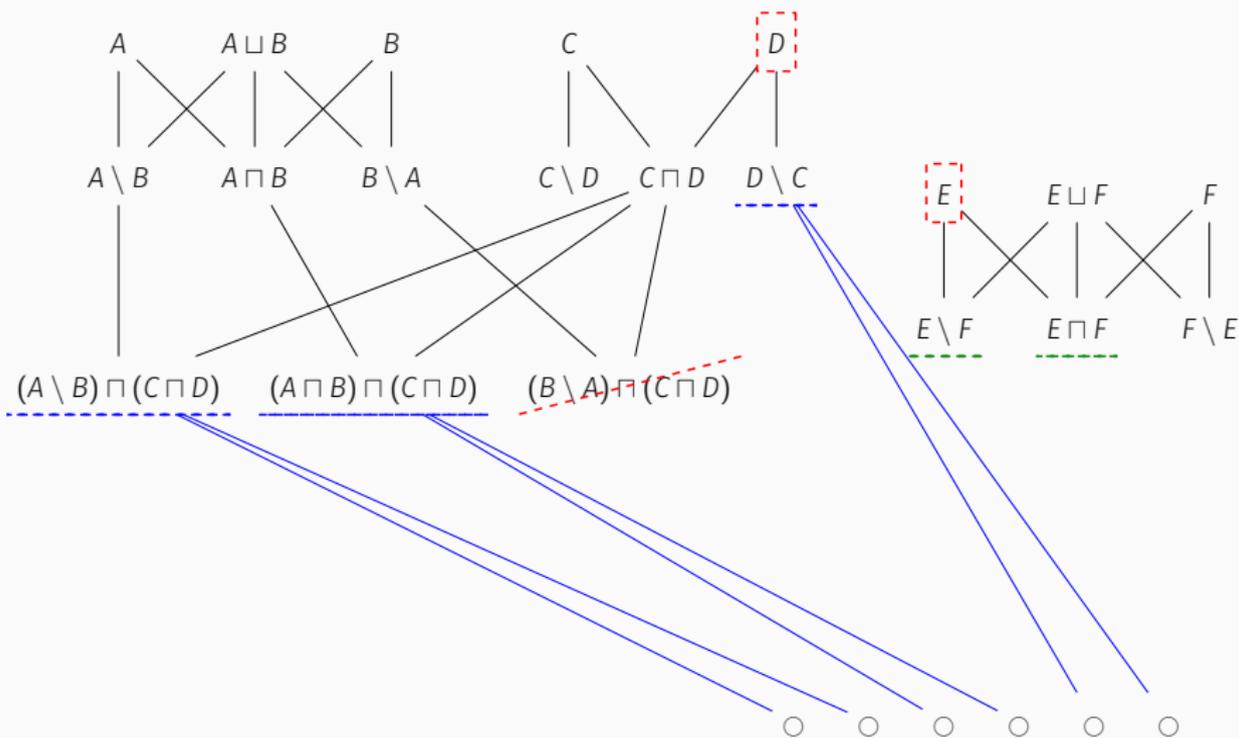
$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 9, |E \sqcup F| \geq 100$



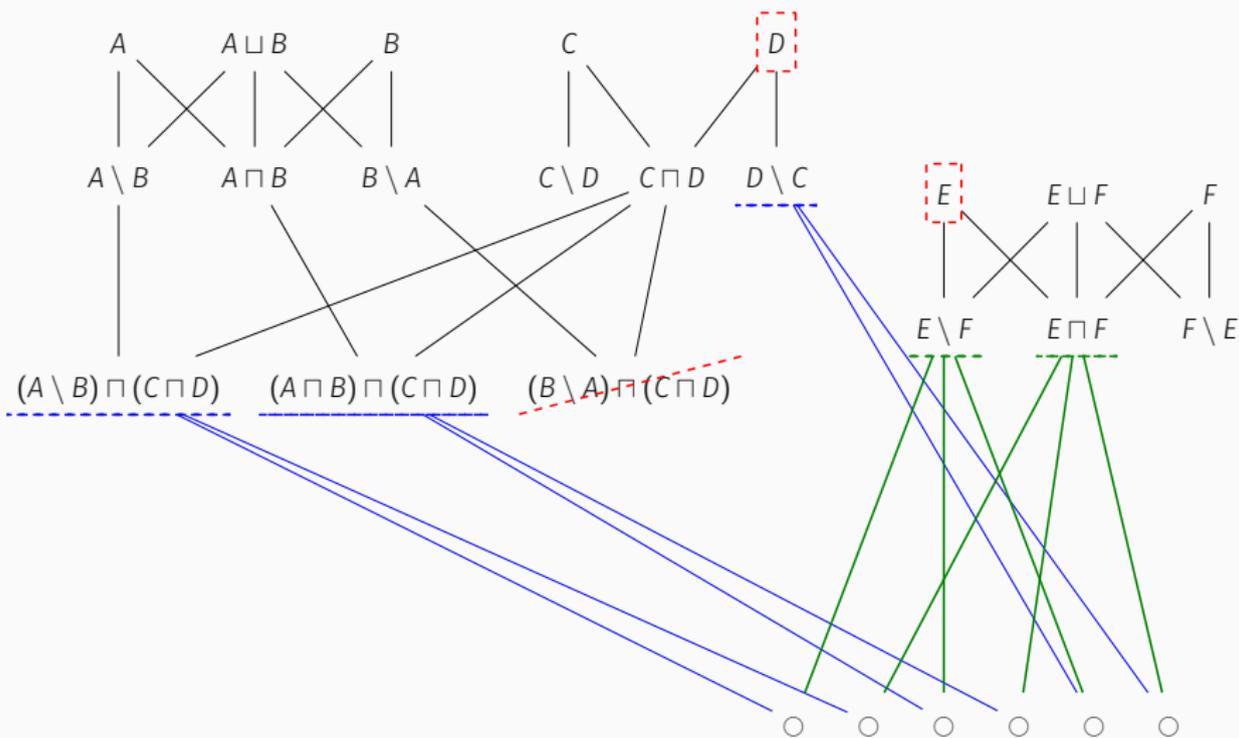
$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 9, |E \sqcup F| \geq 100$



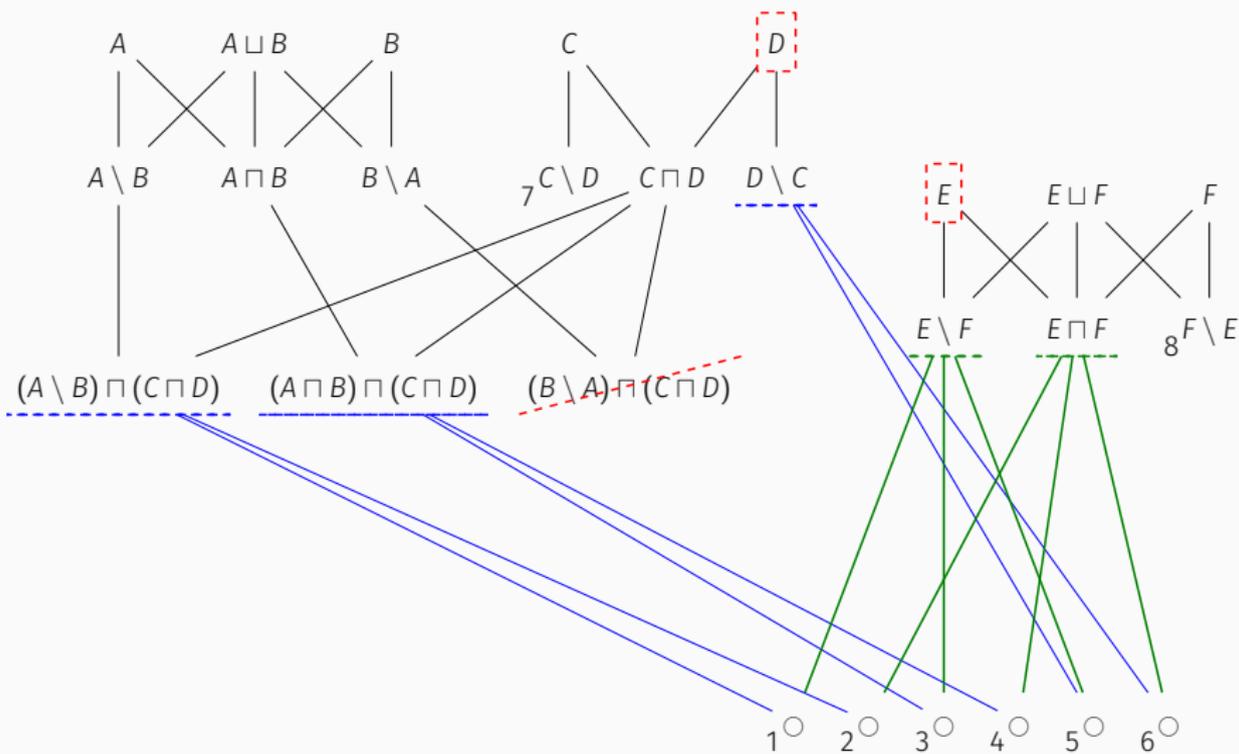
$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \sqcap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 9, |E \sqcup F| \geq 100$



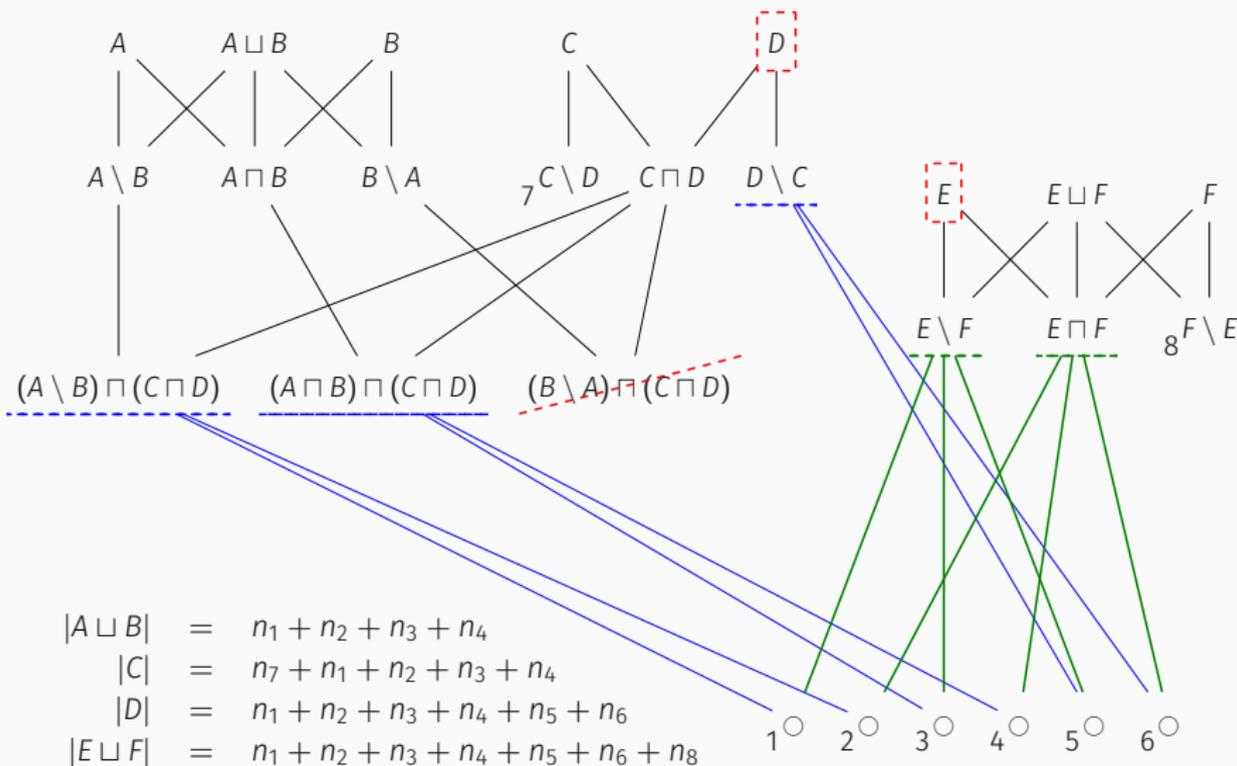
$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 9, |E \sqcup F| \geq 100$



$x \notin D, y \notin A \sqcup B, x \in E \sqcup F, x \in C, y \in D$

$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 9, |E \sqcup F| \geq 100$



Reasoning about cardinality is reduced to solving:

$$\begin{cases} n_1 + n_2 + n_3 + n_4 \geq 4 \\ n_7 + n_1 + n_2 + n_3 + n_4 + n_1 + n_2 + n_3 + n_4 + n_5 + n_6 \leq 10 \\ n_1 + n_2 + n_3 + n_4 + n_5 + n_6 + n_8 \geq 100 \\ n_i \geq 0 \quad \text{for } i \in \{1, \dots, 8\} \end{cases}$$

(which is satisfiable)

Reasoning about cardinality is reduced to solving:

$$\begin{cases} n_1 + n_2 + n_3 + n_4 \geq 4 \\ n_7 + n_1 + n_2 + n_3 + n_4 + n_1 + n_2 + n_3 + n_4 + n_5 + n_6 \leq 10 \\ n_1 + n_2 + n_3 + n_4 + n_5 + n_6 + n_8 \geq 100 \\ n_i \geq 0 \quad \text{for } i \in \{1, \dots, 8\} \end{cases}$$

(which is satisfiable)

However, we must also consider **lower bounds** on region **cardinalities** imposed by **membership** constraints:

$$x \in C, x \notin D \quad \models \quad x \in C \setminus D \quad \models \quad n_7 \geq 1$$

...

Then, arithmetic reasoning can give:

$$n_1 = n_2 = 2 \quad n_3 = n_4 = n_6 = 0 \quad n_5 = n_7 = 1 \quad n_8 = 100$$

Set reasoning can give:

$$r_5 = \{y\} \quad r_7 = \{x\} \quad r_8 = \{x\}$$

Then, arithmetic reasoning can give:

$$n_1 = n_2 = 2 \quad n_3 = n_4 = n_6 = 0 \quad n_5 = n_7 = 1 \quad n_8 = 100$$

Set reasoning can give:

$$r_5 = \{y\} \quad r_7 = \{x\} \quad r_8 = \{x\}$$

We can just add elements to the other regions to match cardinalities:

$$\begin{array}{ll} r_1 = \{e_1, e_2\} & r_2 = \{e_3, e_4\} \\ r_3 = \emptyset & r_4 = \emptyset \\ r_5 = \{e_y\} & r_6 = \emptyset \\ r_7 = \{e_x\} & r_8 = \{e_x, e_5, \dots, e_{103}\} \end{array}$$

Satisfying assignment:

$$\begin{aligned}A &= r_1 \cup r_2 \cup r_3 \cup r_4 &= \{e_1, e_2, e_3, e_4\} \\B &= r_3 \cup r_4 &= \emptyset \\C &= r_1 \cup r_2 \cup r_3 \cup r_4 \cup r_7 &= \{e_x, e_1, e_2, e_3, e_4\} \\D &= r_1 \cup r_2 \cup r_3 \cup r_4 \cup r_5 \cup r_6 &= \{e_y, e_1, e_2, e_3, e_4\} \\E &= r_1 \cup r_2 \cup r_3 \cup r_4 \cup r_5 \cup r_6 &= \{e_y, e_1, e_2, e_3, e_4\} \\F &= r_2 \cup r_4 \cup r_6 \cup r_8 &= \{e_x, e_5, \dots, e_{103}\}\end{aligned}$$

Original problem:

$$x \notin D, y \notin A \sqcup B \quad x \in E \sqcup F, x \in C, y \in D,$$

$$A \sqcup B \approx C \cap D, D \approx E, |A \sqcup B| \geq 4, |C| + |D| \leq 10, |E \sqcup F| \geq 100$$

Our calculus is

- terminating
- refutation sound & complete
- solution sound & complete

for **any** derivation strategy

EXPERIMENTAL EVALUATION 1

file	output	time (s.)	# N	# L
vc1	unsat	0.00	3	3
vc2a	unsat	0.01	17	8
vc2b	sat	0.01	15	7
vc2	unsat	0.00	8	5
vc3a	unsat	0.00	6	0
vc3b	sat	0.01	17	8
vc3	unsat	0.00	6	0
vc4b	sat	0.22	45	16
vc4	unsat	0.07	57	18
vc5b	sat	1.71	71	22
vc5	unsat	0.36	68	21
vc6a	unsat	0.02	34	14
vc6b	sat	0.14	31	13
vc6c	sat	0.06	34	14
vc6	sat	0.02	38	18

- single query benchmarks
- generated by the Jahob system
- from verifying programs with pointer-based data structures

file	output	time (s.)	# N	# L
vc1	1 sat/4 unsat	0.02	12	6
vc2	1 sat/3 unsat	0.07	39	23
vc3	2 sat/2 unsat	0.09	54	21
vc4	1 sat/3 unsat	0.02	0	0
vc5	2 sat/2 unsat	0.08	27	13
vc6	1 sat/3 unsat	0.01	0	0
vc7	2 sat/4 unsat	0.34	56	33
vc8	1 sat/3 unsat	0.01	0	0
vc9	2 sat/2 unsat	0.09	39	19
vc10	2 sat/2 unsat	0.32	94	32

- incremental (i.e., multiple-query) verification conditions
- generated by Leon verification system
- from Scala programs
- mixed constrains over sets, datatypes and bitvectors
- only CVC4 can handle this combination of theories

THANK YOU